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q -Recursive Sequences

Clemens Heuberger

University of Klagenfurt



UNIVERSITÄT
KLAGENFURT

Prelude: Binary Sum of Digits

- Binary expansion $n = \sum_{j \geq 0} \varepsilon_j 2^j$
- Sum of digits $s(n) = \sum_{j \geq 0} \varepsilon_j$

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for $n \geq 0$ with $s(0) = 0$.

In other words:

$$\begin{aligned} s(2n) &= s(n), \\ s(2n + 1) &= s(n) + 1 \end{aligned}$$

for $n \geq 0$ with $s(0) = 0$.

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- Partition into two sets of (almost) equal size;
- Sort both parts individually & recursively;
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Number $M(n)$ of comparisons when sorting n elements:

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for $n \geq 1$ with $M(0) = M(1) = 0$.

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for $n \geq 1$ with $M(0) = M(1) = 0$.

In other words:

$$\begin{aligned}M(2n) &= 2M(n) + 2n - 1, \\M(2n + 1) &= M(n) + M(n + 1) + 2n\end{aligned}$$

for $n \geq 1$ with $M(0) = M(1) = 0$.

(Un-)bordered Factors

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- word w **bordered**:
 - exists non-empty word $v \neq w$
 - v is prefix and suffix of w
- otherwise **unbordered**

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bordered factor	border	length
00	0	2
11	1	2
010	0	3
101	1	3
1010	10	4
0110100110	0110	10

unbordered factor	length
ε	0
0	1
1	1
01	2
10	2
011	3
110	3
100	3
001	3

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bordered factor	border	length
$t[5..6] = 00$	0	2
$t[1..2] = 11$	1	2
$t[3..5] = 010$	0	3
$t[2..4] = 101$	1	3
$t[2..5] = 1010$	10	4
$t[0..9] = 0110100110$	0110	10

unbordered factor	length
ε	0
$t[0..0] = 0$	1
$t[1..1] = 1$	1
$t[0..1] = 01$	2
$t[2..3] = 10$	2
$t[0..2] = 011$	3
$t[1..3] = 110$	3
$t[4..6] = 100$	3
$t[5..7] = 001$	3

Thue–Morse Sequence

$t = 01101001 10010110 10010110 01101001 \dots$

Number of Unbordered Factors

Theorem (Goč–Henshall–Shallit 2013)

*exists unbordered factor
of length n
in Thue–Morse sequence*

\iff

$$(n)_2 \notin 1(01^*0)^*10^*1$$

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- number $f(n)$ of unbordered factors of length n in the Thue–Morse sequence

n	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
$f(n)$	1	2	2	4	2	4	6	0	4	4	4	4	12	0	4	4

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Theorem (Goč–Mousavi–Shallit 2013)

- *inequality $f(n) \leq n$ holds for all $n \geq 4$*
- *$f(n) = n$ infinitely often*
- $\limsup_{n \geq 1} \frac{f(n)}{n} = 1$

Recurrence Relations

- number $f(n)$ of unbordered factors of length n in Thue–Morse sequence
- recurrence relations

$$f(4n) = 2f(2n) \quad (n \geq 2)$$

$$f(4n + 1) = f(2n + 1) \quad (n \geq 0)$$

$$f(8n + 2) = f(2n + 1) + f(4n + 3) \quad (n \geq 1)$$

$$f(8n + 3) = -f(2n + 1) + f(4n + 2) \quad (n \geq 2)$$

$$f(8n + 6) = -f(2n + 1) + f(4n + 2) + f(4n + 3) \quad (n \geq 2)$$

$$f(8n + 7) = 2f(2n + 1) + f(4n + 3) \quad (n \geq 3)$$

Theorem (Goč–Mousavi–Shallit 2013)

$f(n)$ satisfies recurrence relations above

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- $f(n)$ is a 2-recursive sequence

q -Recursive Sequences

Sequence x with

$$x(q^M n + r) = \sum_{\ell \leq k \leq u} c_{r,k} x(q^m n + k)$$

for all $n \geq n_0$ and $0 \leq r < q^M$.

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Constants:

- $q \geq 2$, $M > m \geq 0$, $\ell \leq u$, $n_0 \geq \max\{-\ell/q^m, 0\}$ integers
- $c_{s,k} \in \mathbb{C}$ for all $0 \leq s < q^M$ and $\ell \leq k \leq u$

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q -recursive sequence

q -Regular Sequences

Vector-valued sequence $v: \mathbb{N}_0 \rightarrow \mathbb{C}^D$ with

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Every q -recursive sequence is q -regular.

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Regular sequences as inhomogeneities in the recurrence of the q -recursive sequence are allowed.

Minimality of Linear Representations

- Given regular sequence x , find matrices A_0, \dots, A_r , row vector u , column vector w and vector valued sequence v of minimal dimension D such that $x(n) = uv(n)$, $v(0) = w$, and $v(qn + r) = A_r v(n)$ for all $n \geq 0$ and $0 \leq r < q$.

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- These conditions are also **sufficient** (Schützenberger 1961; Berstel–Reutenauer 2011; H.–Krenn–Lipnik 202?).

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- Implemented in SageMath.